Using a Linux Security Module for contest security

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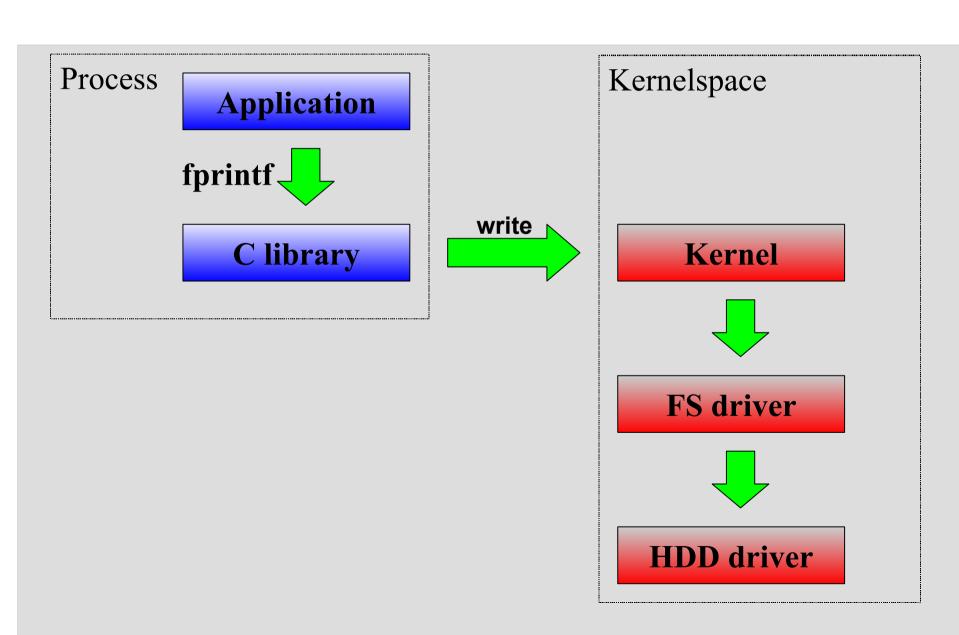
- Goals
- Background
- Overview of techniques
- Our implementation
- Java
- Conclusions

The goals

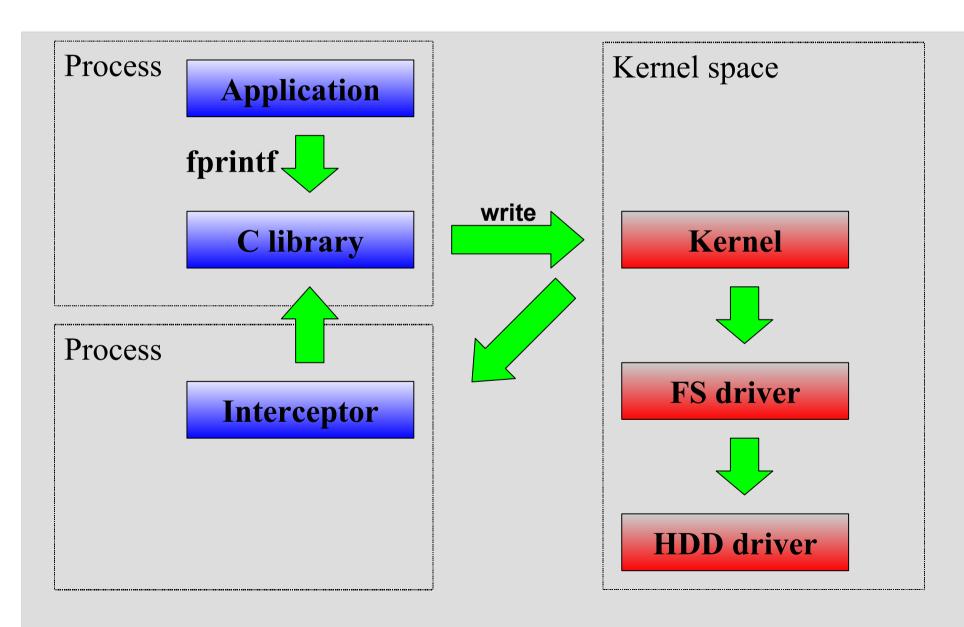
- Resource limits
- No networking
- No IPC
- No access to evaluation system
- Single process
- Single thread

- Accurate constraints
- High throughput
- Minimum overhead
- Transparent

Device access in Linux



System call wrappers



System call wrappers

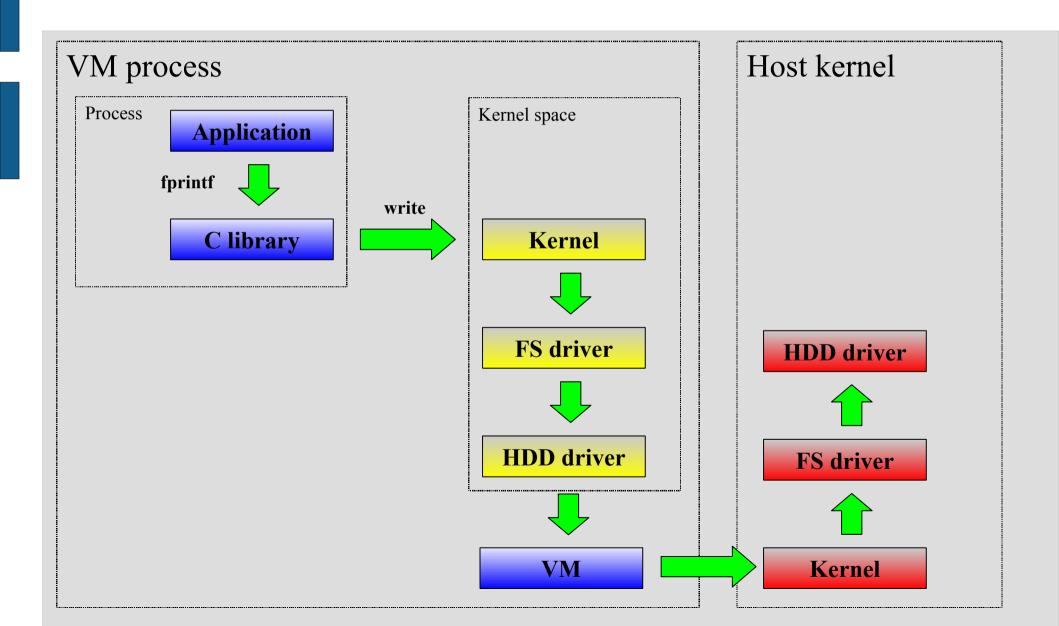
Pros

- Configurable, offthe-shelf wrappers available
- Minimal startup overhead

Cons

- Context switch per system call
- Huge number of system calls
- Poor security track record

Virtualisation



Virtualisation

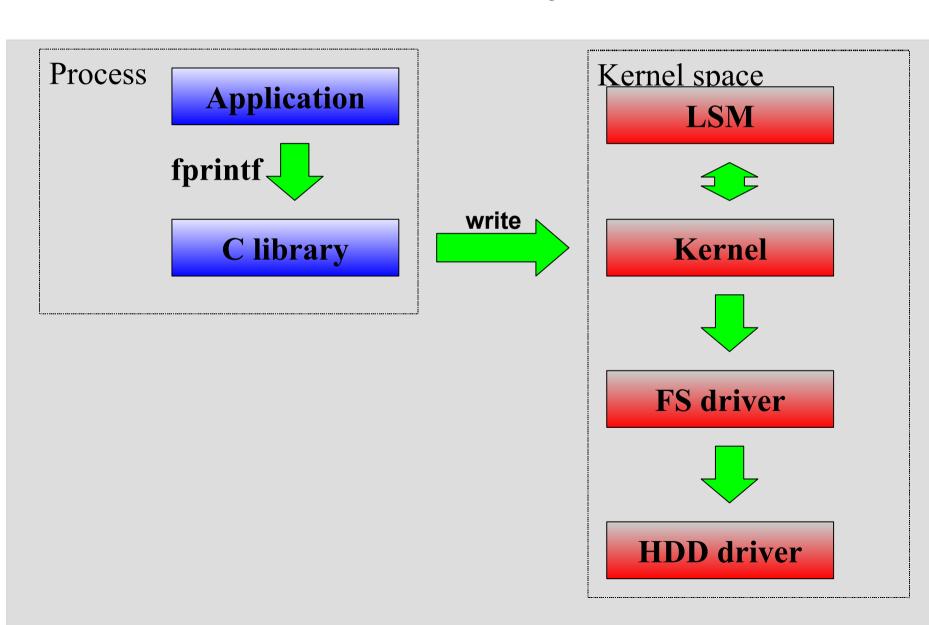
Pros

- Guest OS can be totally isolated
- Can start with a totally fresh OS for each run

Cons

- Performance impact
- Startup time
- Does not prevent multi-threading, external processes etc.

Linux Security Module



Linux Security Module

Pros

Policy per operation, not per syscall

- No extra context switches
- Access to kernel internals
- Fewer races

Cons

- Kernel programming is difficult
- Interface changes frequently
- Outdated docs

LSM implementation: sandtray

- Launcher program requests restrictions:
 - Calls setrlimit to set CPU, memory etc limits
 - Writes to /proc/self/attr/exec to set further limits:
 - version 1.0 (sets default restrictions)
 - allow write problem.out
- Launcher then calls exec
 - This triggers sandtray for this process
- Caller asks for exact CPU time on return
 - setrlimit only has 1 second resolution

Filesystem access

- glibc accesses huge numbers of files
 - A whitelist is difficult to maintain
 - Path-based checks tricky due to links
- Instead, read access left open
 - Contest internals owned by a different user
- Write access is tightly controlled
 - Only the output file may be written
 - Together with setrlimit, limits total disk space
 - No symlinks, chmod, chown, etc.

Covert channels

- Sandtray cuts off
 - Networking, SystemV IPC, kill etc.
 - Writing files into / tmp or similar
- Are still channels through /proc and others
 - Now prevented by serialising execution
- Cache timing theoretically possible
 - Probably harder than solving the original problem

Java (Sun VM)

- Consumes huge amounts of virtual memory
- Does lots of suspicious-looking things
- Has its own security manager
 - permission java.io.FilePermission ...
- Has command-line option for max heap size:
 - -Xm \times 64m
- We use these instead of sandtray

Conclusions

LSM

- Good abstraction of operations
- Low overhead
- Interface is a moving target
- Sun Java VM
 - VM does not play nicely with LSM
 - Internal security tools are good enough

Questions?